

Generalization of proofs and codification of graph families

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Contents

1 Generalization of proofs

2 Codification of graph families

Fermat numbers

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- Only $F(0), \dots, F(11)$ are fully factored.
- $F(20)$ and $F(24)$ are known to be composite, but no prime factor is known.

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$$a_1 \mid b_1 \cdot d_1 + c_1 \cdot d_1, a_2 \mid b_2 + c_2$$

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$$\begin{array}{l} 641 \mid (5^4 \cdot 2^{28} - 1) + (2^{32} + 1) \\ a_2 \mid (b_2 - d_2) + (c_2 + d_2), a_3 \mid b_3 + c_3 \\ A \mid (D \cdot C - E) + (B \cdot C + E) \end{array}$$

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Application to Fermat numbers

- Original calculation: $\left\{ \begin{array}{l} 641 \mid 5^4 + 2^4 \\ 641 \mid 5^4 \cdot 2^{28} - 1 \end{array} \right. \Rightarrow 641 \mid 2^{32} + 1.$

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- Particularization: $k \cdot 2^s + 1 \mid k^{2 \cdot r} + 2^{2^n - 2 \cdot r \cdot s} \Rightarrow k \cdot 2^s + 1 \mid F(n)$
 $(A = k \cdot 2^s + 1, B = 2^{2^n - 2 \cdot r \cdot s}, C = 2^{2 \cdot r \cdot s}, D = k^{2 \cdot r}, E = 1).$
 Note that, by varying r , we can make $k^{2 \cdot r} + 2^{2^n - 2 \cdot r \cdot s}$ way smaller than $F(n)$.

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 Note that, by varying r , we can make $k^{2 \cdot r} + 2^{2^n - 2 \cdot r \cdot s}$ way smaller than $F(n)$.
- Remark: this method can be applied in any area of mathematics.

Other results

- Result 1:

$$\left\{ \begin{array}{l} d \mid k \cdot 2^s \\ (k \cdot 2^s + 1) \cdot d^{2 \cdot t} - (k \cdot 2^s)^{2 \cdot t} + 1 = F(n) \end{array} \Rightarrow k \cdot 2^s + 1 \mid F(n) \right.$$

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- Result 2:
$$(k \cdot 2^s + 1 - k^{2 \cdot q}) \cdot 2^{2 \cdot q \cdot s} + 1 = F(n) \Rightarrow k \cdot 2^s + 1 \mid F(n).$$

This theorem was obtained by Baaz in 1999.

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- Result 3:
$$p \mid (F(n-1) - 1 - p \cdot c)^2 + 1 \Rightarrow p \mid F(n).$$

Note that, by varying c , we can make $(F(n-1) - 1 - p \cdot c)^2 + 1$ way smaller than $F(n)$.

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Note that, by varying c , we can make $(F(n-1) - 1 - p \cdot c)^2 + 1$ way smaller than $F(n)$.

- Result 4:
$$p \nmid \frac{(F(n+2) - 1)^k - 1}{F(n)} \Leftrightarrow p \mid F(n),$$

where p equals $k \cdot 2^{n+2} + 1$ and is prime.

The converse implication was proved by Wang.

Contents

1 Generalization of proofs

2 Codification of graph families

Vertex-by-vertex increasing graph sequences

- A sequence G of simple graphs is vertex-by-vertex increasing if and only if $V(G_n) = \{1, \dots, n+1\}$ and $E(G_n) \subseteq E(G_{n+1})$.

The bijection

Input: vertex-by-vertex increasing graph sequence G such that G_∞ is Collatz's graph.

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G_1	$\text{adj}(G_1)$	
$1 — 2$	$\begin{bmatrix} * & * \\ 1 & * \end{bmatrix}$	$(\Phi(G))_0 = 1$

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G_2	$\text{adj}(G_2)$	
1 —— 2 3	$\begin{bmatrix} * & * & * \\ 1 & * & * \\ 0 & 0 & * \end{bmatrix}$	$(\Phi(G))_0 = 1$ $(\Phi(G))_1 = 0$

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G_3	$\text{adj}(G_3)$	
 $\begin{array}{c} 1 \quad 2 \\ \diagdown \quad \diagup \\ \qquad \\ 4 \end{array}$	$\begin{bmatrix} * & * & * & * \\ 1 & * & * & * \\ 0 & 0 & * & * \\ 1 & 1 & 0 & * \end{bmatrix}$	$(\Phi(G))_0 = 1$ $(\Phi(G))_1 = 0$ $(\Phi(G))_2 = 6$

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G_4	$\text{adj}(G_4)$	
 $\begin{array}{ccccc} 1 & \xrightarrow{\hspace{1cm}} & 2 & & 3 \\ & & \diagup & & \\ 4 & & 5 & & \end{array}$	$\begin{bmatrix} * & * & * & * & * \\ 1 & * & * & * & * \\ 0 & 0 & * & * & * \\ 1 & 1 & 0 & * & * \\ 0 & 0 & 0 & 0 & * \end{bmatrix}$	$(\Phi(G))_0 = 1$ $(\Phi(G))_1 = 0$ $(\Phi(G))_2 = 6$ $(\Phi(G))_3 = 0$

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G_5	$\text{adj}(G_5)$	
<pre> graph LR 1 --- 2 1 --- 4 2 --- 3 4 --- 5 5 --- 6 </pre>	$\begin{bmatrix} * & * & * & * & * & * \\ 1 & * & * & * & * & * \\ 0 & 0 & * & * & * & * \\ 1 & 1 & 0 & * & * & * \\ 0 & 0 & 0 & 0 & * & * \\ 0 & 0 & 1 & 0 & 0 & * \end{bmatrix}$	$(\Phi(G))_0 = 1$ $(\Phi(G))_1 = 0$ $(\Phi(G))_2 = 6$ $(\Phi(G))_3 = 0$ $(\Phi(G))_4 = 4$

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G_6	$\text{adj}(G_6)$	
<pre> graph LR 1 --- 2 1 --- 4 2 --- 3 4 --- 5 5 --- 6 </pre>	$\left[\begin{array}{ccccccc} * & * & * & * & * & * & * \\ 1 & * & * & * & * & * & * \\ 0 & 0 & * & * & * & * & * \\ 1 & 1 & 0 & * & * & * & * \\ 0 & 0 & 0 & 0 & * & * & * \\ 0 & 0 & 1 & 0 & 0 & * & * \\ 0 & 0 & 0 & 0 & 0 & 0 & * \end{array} \right]$	$(\Phi(G))_0 = 1$ $(\Phi(G))_1 = 0$ $(\Phi(G))_2 = 6$ $(\Phi(G))_3 = 0$ $(\Phi(G))_4 = 4$ $(\Phi(G))_5 = 0$
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G_7	$\text{adj}(G_7)$	
<pre> graph LR 1 --- 2 1 --- 4 2 --- 3 4 --- 5 5 --- 6 5 --- 4 6 --- 3 8 --- 7 </pre>	$\text{adj}(G_7) = \begin{bmatrix} * & * & * & * & * & * & * & * \\ 1 & * & * & * & * & * & * & * \\ 0 & 0 & * & * & * & * & * & * \\ 1 & 1 & 0 & * & * & * & * & * \\ 0 & 0 & 0 & 0 & * & * & * & * \\ 0 & 0 & 1 & 0 & 0 & * & * & * \\ 0 & 0 & 0 & 0 & 0 & 0 & * & * \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & * \end{bmatrix}$	$(\Phi(G))_0 = 1$ $(\Phi(G))_1 = 0$ $(\Phi(G))_2 = 6$ $(\Phi(G))_3 = 0$ $(\Phi(G))_4 = 4$ $(\Phi(G))_5 = 0$ $(\Phi(G))_6 = 8$

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G_8	$\text{adj}(G_8)$	
<pre> graph LR 1 --- 2 1 --- 5 2 --- 3 4 --- 5 5 --- 6 7 --- 8 8 --- 9 </pre>	$\begin{bmatrix} * & * & * & * & * & * & * & * & * \\ 1 & * & * & * & * & * & * & * & * \\ 0 & 0 & * & * & * & * & * & * & * \\ 1 & 1 & 0 & * & * & * & * & * & * \\ 0 & 0 & 0 & 0 & * & * & * & * & * \\ 0 & 0 & 1 & 0 & 0 & * & * & * & * \\ 0 & 0 & 0 & 0 & 0 & 0 & * & * & * \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & * & * \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & * \end{bmatrix}$	$(\Phi(G))_0 = 1$ $(\Phi(G))_1 = 0$ $(\Phi(G))_2 = 6$ $(\Phi(G))_3 = 0$ $(\Phi(G))_4 = 4$ $(\Phi(G))_5 = 0$ $(\Phi(G))_6 = 8$ $(\Phi(G))_7 = 0$

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G_9	$\text{adj}(G_9)$	
<pre> graph LR 1 --- 2 1 --> 4 2 --- 3 2 --> 5 4 --- 5 4 --- 8 5 --- 6 5 --> 10 6 --- 10 7 --- 9 8 --- 10 </pre>	$\begin{bmatrix} * & * & * & * & * & * & * & * & * & * \\ 1 & * & * & * & * & * & * & * & * & * \\ 0 & 0 & * & * & * & * & * & * & * & * \\ 1 & 1 & 0 & * & * & * & * & * & * & * \\ 0 & 0 & 0 & 0 & * & * & * & * & * & * \\ 0 & 0 & 1 & 0 & 0 & * & * & * & * & * \\ 0 & 0 & 0 & 0 & 0 & 0 & * & * & * & * \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & * & * & * \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & * & * \\ 0 & 0 & 1 & 0 & 1 & 0 & 0 & 0 & 0 & * \end{bmatrix}$	$(\Phi(G))_0 = 1$ $(\Phi(G))_1 = 0$ $(\Phi(G))_2 = 6$ $(\Phi(G))_3 = 0$ $(\Phi(G))_4 = 4$ $(\Phi(G))_5 = 0$ $(\Phi(G))_6 = 8$ $(\Phi(G))_7 = 0$ $(\Phi(G))_8 = 80$

Application to Collatz Problem

- Output: $(\Phi(G))_n = \begin{cases} 0, & \text{if } n \text{ is odd;} \\ \sqrt{2^n} + \sqrt[3]{4^{n+1}}, & \text{if } n \in \{6 \cdot r + 2\}_{r \in \mathbb{N}}; \\ \sqrt{2^n}, & \text{otherwise,} \end{cases}$
(first terms: 1, 0, 6, 0, 4, 0, 8, 0, 80...).

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- Collatz Problem: is G_∞ connected?
- Pattern observed: if $n \geq 3$, then the number of connected components of G_n is apparently
 $\left\lfloor \frac{n+2}{2} \right\rfloor - \left\lfloor \frac{n+3}{6} \right\rfloor + 1.$

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 $\left\lfloor \frac{n+2}{2} \right\rfloor - \left\lfloor \frac{n+3}{6} \right\rfloor + 1.$
- Remark: because Φ is a bijection, we can input integer sequences instead
 (provided that they are upper-bounded by 2^{n+1}).