



NATIONAL RESEARCH
UNIVERSITY

July 13, 2021

SHIFTING ANY PATH TO AN AVOIDABLE ONE

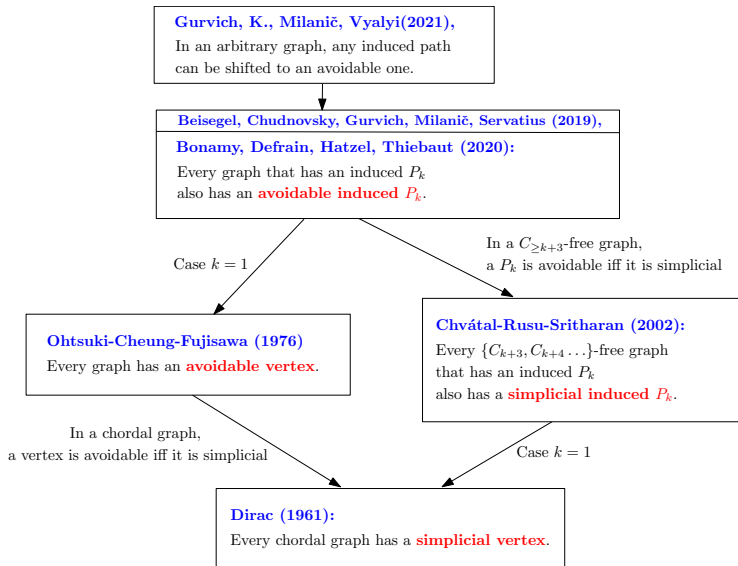
Matjaž Krnc
University of Primorska

joint work with V. Gurvich,
M. Milanič and M. Vyalyi

8th European Congress of
Mathematics

Algorithmic Graph Theory (MS - ID 54)

OUTLINE

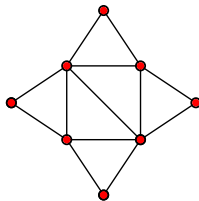
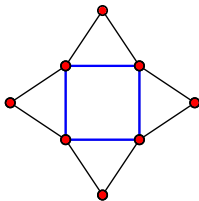


CHORDAL GRAPHS

A graph G is **chordal** if every cycle in G of length at least four has a chord.

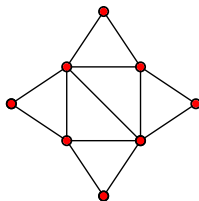
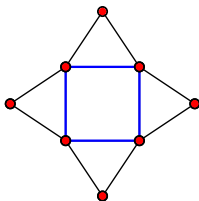
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Chordal graphs are known to have several structural and algorithmic properties.

CHORDAL GRAPHS

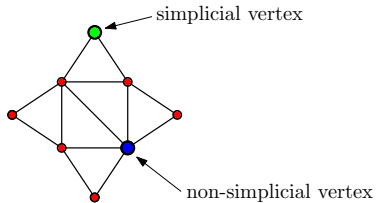
Theorem (Dirac 1961)

*Every chordal graph has a **simplicial vertex**, that is, a vertex whose neighborhood is a clique.*

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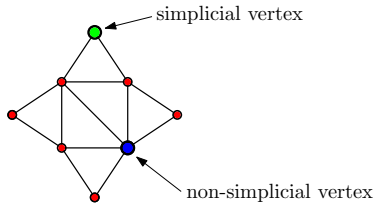
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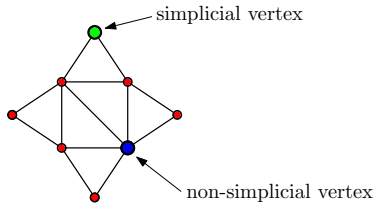


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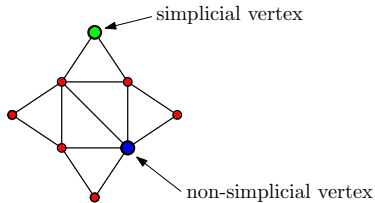
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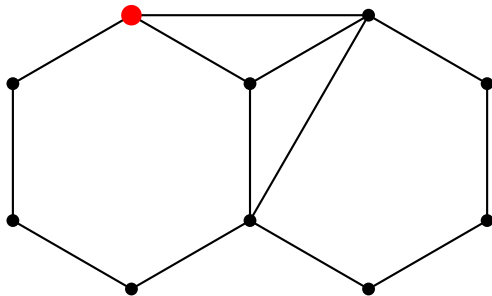
Simplicial vertices provide another characterization of chordal graphs, and were generalized in the literature in various ways:

- by generalizing **simpliciality** to a concept in general graphs;
- by generalizing the 'simpliciality' property from vertices, which are paths of length 0, to **longer induced paths**.

First generalization: from chordal graphs to all graphs

AVOIDABLE VERTICES

This vertex is avoidable:



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In a **tree** ($\neq K_1$) a vertex is avoidable if and only if it is a **leaf**.

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Theorem (Ohtsuki, Cheung, Fujisawa, 1976)

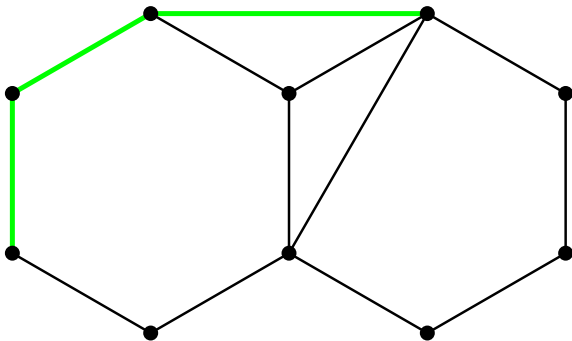
Every graph has an avoidable vertex.

Second generalization: from vertices to longer paths

SIMPLICIAL PATHS

Definition

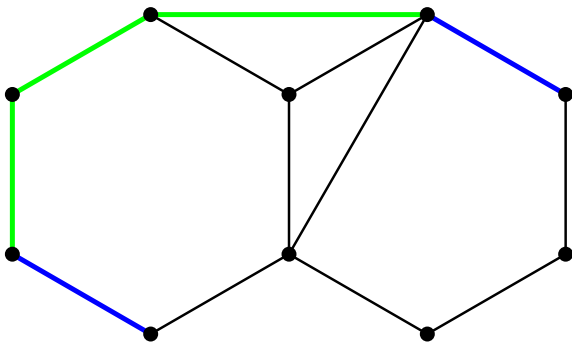
Given an induced path P in a graph G , an **extension** of P is any induced path in G obtained by adding to P one edge at each end.



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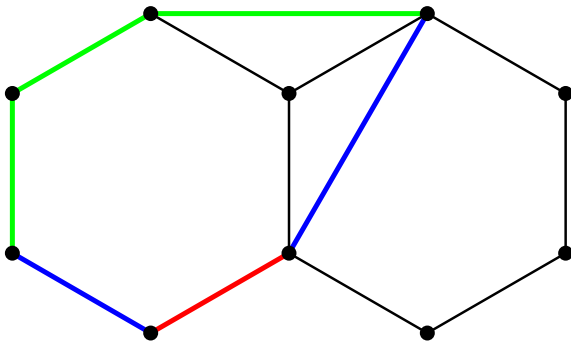
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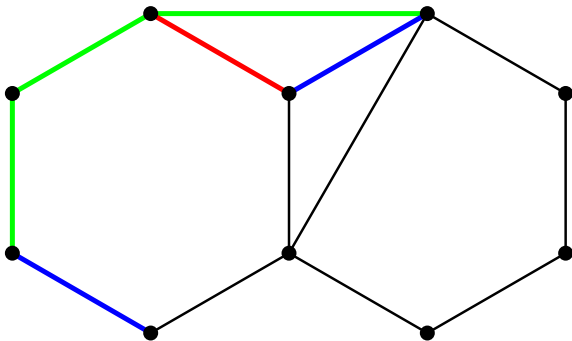
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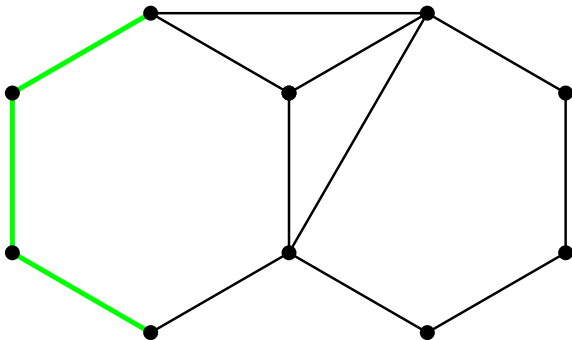


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Given an induced path P in a graph G , an **extension** of P is any induced path in G obtained by adding to P one edge at each end.

This green path is simplicial:

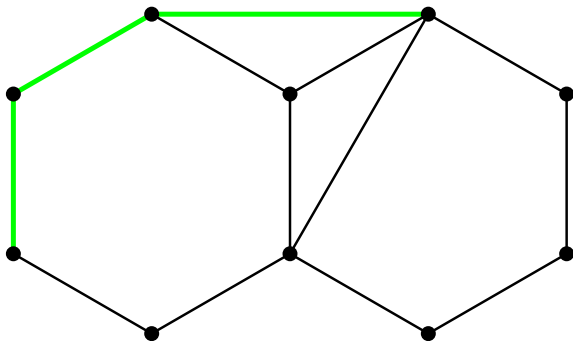


SIMPLICIAL PATHS

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Given an induced path P in a graph G , an **extension** of P is any induced path in G obtained by adding to P one edge at each end.

However, this other green path is not simplicial:



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Theorem (Chvátal, Rusu, Sritharan, 2002)

For each $k \geq 1$, every graph without induced cycles of length more than $k + 2$ that has an induced P_k , has a simplicial induced P_k .

GENERALIZATIONS OF SIMPLICIAL VERTICES

Ohtsuki-Cheung-Fujisawa (1976)

Every graph has an **avoidable vertex**.

In a chordal graph,
a vertex is avoidable iff it is simplicial

Chvátal-Rusu-Sritharan (2002):

Every $\{C_{k+3}, C_{k+4}, \dots\}$ -free graph
that has an induced P_k
also has a **simplicial induced P_k** .

Case $k = 1$

Dirac (1961):

Every chordal graph has a **simplicial vertex**.

**A common generalization:
avoidable paths!**

AVOIDABLE PATHS

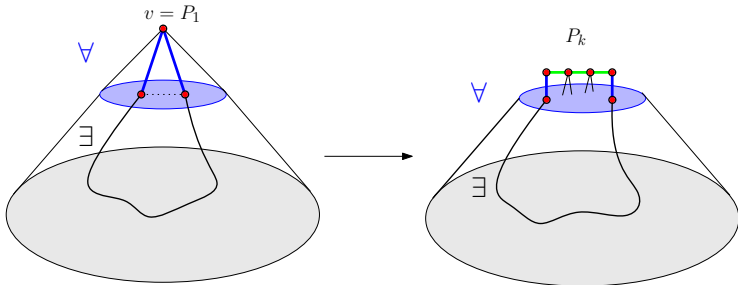
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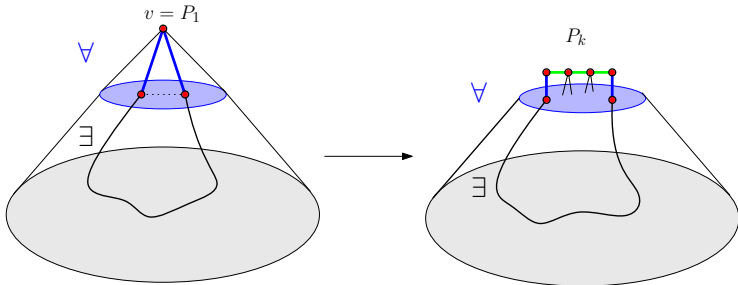
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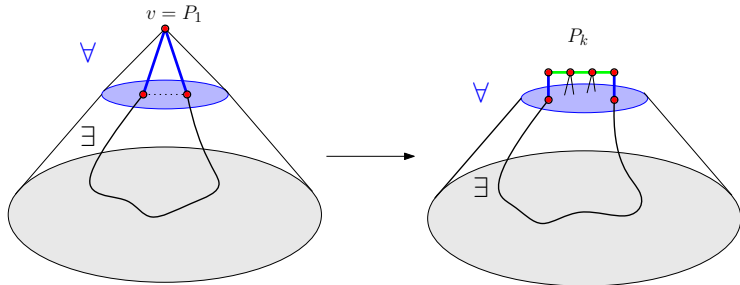


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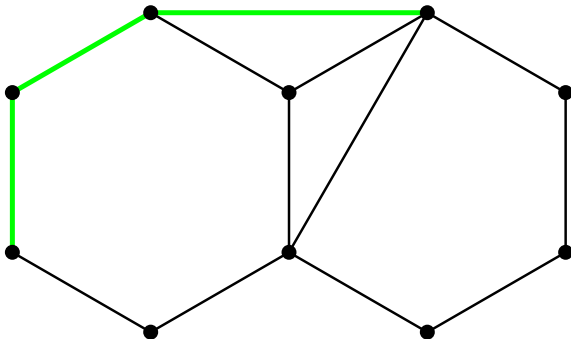


avoidable vertex $v \equiv$ avoidable P_1

Every simplicial induced path is (trivially) avoidable.

AVOIDABLE PATHS

This green path is not avoidable:



AVOIDABLE PATHS

Conjecture (Beisegel, Chudnovsky, Gurvich, Milanič, Servatius)

For every $k \geq 1$, every graph that has an induced P_k also has an avoidable induced P_k .

Theorem (Beisegel, Chudnovsky, Gurvich, Milanič, Servatius)

Every graph that has an edge also has an avoidable edge.

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Theorem (Bonamy, Defrain, Hatzel, Thiebaut (2020))

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COMMON GENERALIZATION

Beisegel, Chudnovsky, Gurvich, Milanič, Servatius (2019),

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Case $k = 1$

Shifting any path to an avoidable one

AVOIDABLE PATHS AND SHIFTING

Let P and P' be induced k -paths in G .

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Theorem (Gurvich, K., Milanič, Vyalyi)

In every graph, every induced path can be shifted to an avoidable one.

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For a graph G and a positive integer k ,

- property $H_B(G, k)$ holds if every induced P_k in G can be shifted to an avoidable one;

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$H_R(G, k)$ implies $H_B(G, k)$.

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So it is enough to prove:

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Recall the ingredients:

- property $H_B(G, k)$ – every induced P_k in G can be shifted to an avoidable one;
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- property $H_R(G, k)$ – if $H_R(G, k, v)$, for every $v \in V(G)$.
- **Lemma:** $H_R(G, k)$ implies $H_B(G, k)$.
- **Lemma [BDHT (2020)]:** Let $G' = G|_{u_1 u_2 \rightarrow u}$. If an induced k -path Q in $G' - N[u]$ is avoidable in G' , it is avoidable in G as well.

Procedure 1 SHIFTING(G, P)

Input: a graph G and an induced path $P = p_1p_2 \dots p_k$ in G

Output: a sequence S of paths shifting P to an avoidable induced path in G

- 1: **if** there exists an extension xPy of P **then**
 - 2: $Q \leftarrow yp_k \dots p_1x$
 - 3: **return** $P, \text{REFINEDSHIFTING}(G, Q)$
 - 4: **else**
 - 5: **return** P
-

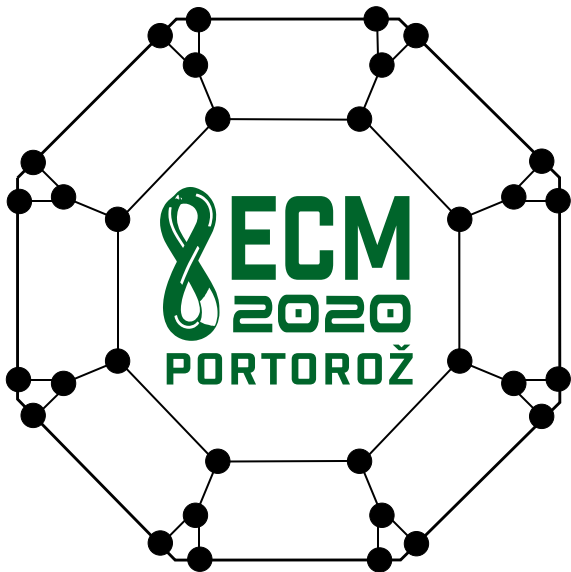
Procedure 2 REFINEDSHIFTING(G, P)

Input: a graph G and an induced path $P = p_1 \dots p_{k+2}$ in G

Output: a sequence S of paths in $G - N[p_{k+2}]$ shifting $p_1 \dots p_k$ to an avoidable induced path in G

- 1: $P' \leftarrow p_1 \dots p_k$
 - 2: $S \leftarrow$ the one-element sequence containing path P'
 - 3: **if** there exists an extension $xP'y$ in $G - N[p_{k+2}]$ **then**
 - 4: $S \leftarrow S, \text{REFINEDSHIFTING}(G - N[p_{k+2}], xP'y)$
 - 5: $Q \leftarrow$ the end path of S
 - 6: **if** Q has an extension xQy in G such that y is the unique neighbor of p_{k+2} in $\{x, y\}$ **then**
 - 7: let $Q = q_1 \dots q_k$ such that y is adjacent to q_k
 - 8: $Q' \leftarrow xq_1 \dots q_k$
 - 9: $G' \leftarrow G /_{p_{k+2}y \rightarrow y'}$
 - 10: $S' \leftarrow \text{REFINEDSHIFTING}(G', Q'y')$
 - 11: **return** S, S'
 - 12: **else**
 - 13: **return** S
-

THANK YOU!



Find a shortest induced path which is not avoidable.