### On coverings and perfect colorings of hypergraphs

#### Anna Taranenko

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# Perfect coloring of graphs

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#### Our goal: develop the similar concepts for hypergraphs

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# Hypergraphs, incidence matrices, bipartite representation

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The bipartite representation G(X, W; E) of a hypergraph  $\mathcal{G}(X, W)$  is a bipartite graph, x is adjacent to w in G iff x is incident to w in  $\mathcal{G}$ . The adjacency matrix  $M_G$  of the bipartite representation is

$$M_G = \left(\begin{array}{cc} 0 & \mathbb{I} \\ \mathbb{I}^T & 0 \end{array}\right).$$

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Another approach: the adjacency matrix of a *d*-uniform hypergraph is a *d*-dimensional matrix.

A *d*-dimensional matrix *A* of order *n* is an array  $(a_{\alpha})$ ,  $\alpha = (\alpha_1, \ldots, \alpha_d)$ ,  $\alpha_i \in \{1, \ldots, n\}$ .

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**1. Combinatorial approach:** The adjacency matrix  $\mathbb{M}$  of a *d*-uniform hypergraph  $\mathcal{G}$  on *n* vertices is a *d*-dimensional (0, 1)-matrix of order *n* with entries  $m_{\alpha} = 1 \Leftrightarrow (\alpha_1, \ldots, \alpha_d)$  is a hyperedge in  $\mathcal{G}$ .

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**2. Algebraic approach:** For 2-dimensional matrices  $B^1, \ldots, B^d$  of order *n* define folding  $C = [B^1, \ldots, B^d]$  to be the *d*-dimensional matrix of order *n*:

$$c_{lpha_1,...,lpha_d} = \sum_{i=1}^n b^1_{lpha_1,i}\cdots b^d_{lpha_d,i}.$$

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The full adjacency matrix  $\mathbb{A}$  of a *d*-uniform hypergraph  $\mathcal{G}$  on *n* vertices is

$$\mathbb{A} = [\mathbb{I}, \dots, \mathbb{I}].$$

# Full adjacency matrices and totally regular hypergraphs

Proposition

Let  $\mathbb{A}$  be the full adjacency matrix of a *d*-uniform hypergraph. Then entries  $a_{\alpha}$  are exactly the degrees of sets  $S(\alpha) = \{\alpha_1, \ldots, \alpha_d\}$ .

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We will say that a *d*-uniform hypergraph  $\mathcal{G} = (X, W)$  is totally  $(r_1, \ldots, r_{d-1})$ -regular if every  $S \subset X$ , |S| = i, has  $deg(S) = r_i$ .

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#### Proposition

If  $\mathcal G$  is a d-uniform totally  $(r_1,\ldots,r_{d-1})$ -regular hypergraph, then

$$\mathbb{A}=\mathbb{M}+\sum_{t=1}^{d-1}r_t\mathcal{I}_t,$$

where  $\mathcal{I}_t$  is a *d*-dimensional (0,1)-matrix, whose unity entries indexed by  $(\alpha_1, \ldots, \alpha_d)$  with exactly *t* different components.

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• If v is a vector, then  $A \circ v$  is a vector u:

$$u_j = \sum_{i_1,\dots,i_{d-1}=1}^n a_{j,i_1,\dots,i_{d-1}} v_{i_1} \cdots v_{i_{d-1}}$$

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• If P is a 2-dimensional matrix, then  $A \circ P$  is a d-dimensional matrix C:

$$C_{j,k_1...,k_{d-1}} = \sum_{i_1,...,i_{d-1}=1}^n a_{j,i_1,...,i_{d-1}} p_{k_1,i_1} \cdots p_{k_{d-1},i_{d-1}}.$$

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• If B a t-dimensional matrix of order n, then  $A \circ B$  is a similar ((d-1)(t-1)+1)-dimensional matrix.

 $\lambda$  is an eigenvalue and  $\nu$  is the eigenvector of a *d*-dimensional matrix A if

 $A \circ v = \lambda(\mathcal{I} \circ v),$ 

where  $\mathcal{I} = \mathcal{I}_1$  is the *d*-dimensional identity matrix.

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Theorem (T., 2021+)

Let v be a plain eigenvector for  $\mathcal{G}$ . Then

• v is a full eigenvector for  $\mathcal{G}$ ;

• If  $\mathcal{G}$  is a totally regular hypergraph, then v is an eigenvector for  $\mathcal{G}$ .

In all cases, eigenvalues of one type can be counted from another.

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**Using bipartite representation:** A perfect *k*-coloring of a G is given by vertex *k*-coloring matrix *P* and hyperedge coloring matrix *R* satisfying

$$\left(\begin{array}{cc} 0 & \mathbb{I} \\ \mathbb{I}^T & 0 \end{array}\right) \left(\begin{array}{cc} 0 & P \\ R & 0 \end{array}\right) = \left(\begin{array}{cc} 0 & P \\ R & 0 \end{array}\right) \left(\begin{array}{cc} 0 & S \\ T & 0 \end{array}\right).$$

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All three definitions are equivalent.

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HV- and VH-parameter matrices:

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The parameter matrix S:

$$S_{lpha}=r\prod_{i=2}^{d}t^{lpha_{i}}(d-t)^{1-lpha_{i}}; \hspace{0.2cm} lpha_{i}\in\{0,1\}.$$

### Parameter matrix of a hypergraph

Let  $\mathcal{G}$  be a *d*-uniform hypergraph with the full adjacency matrix  $\mathbb{A}$ . *P* is a perfect *k*-coloring of  $\mathcal{G}$  if

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Theorem (T., 2021+)

S is the HV-parameter matrix and T is the VH-parameter matrix of a perfect coloring of a hypergraph G if and only if

 $\mathbb{S} = [T, S^T, \dots, S^T].$ 

In case of totally regular hypergraphs, perfect colorings can be defined as through the adjacency matrix, as the full adjacency matrix.

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#### Proposition

If  $\mathcal{G}$  is a uniform totally regular hypergraph, P is a vertex coloring matrix, then

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### Theorem (T., 2021+)

Let *P* be a perfect coloring of a hypergraph  $\mathcal{G}$  with the parameter matrix  $\mathbb{S}$ . If  $\lambda$  and v are eigenvalue and eigenvector of  $\mathbb{S}$ , then  $\lambda$  and v are eigenvalue and eigenvector of the full adjacency matrix  $\mathbb{A}$ .

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### Theorem (T., 2021+)

The parameter matrix S of a perfect coloring of a hypergraph can be symmetrized by the vector of color densities.

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Coverings and colorings of hypergraphs

### Theorem (Weisfeiler, Leman, 1968)

Let G be a graph. Then there is the minimal perfect coloring f such that every other perfect coloring of G is obtained from f by splitting of color classes.

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#### Theorem (T., 2021+)

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### Theorem (Weisfeiler, Leman, 1968)

Let G be a graph. Then there is the minimal perfect coloring f such that every other perfect coloring of G is obtained from f by splitting of color classes.

For regular graphs the minimal perfect coloring is monochromatic.

#### Theorem (T., 2021+)

Let  $\mathcal{G}$  be a hypergraph. Then there is the minimal perfect coloring f such that every other perfect coloring of  $\mathcal{G}$  is obtained from f by slitting of color classes.

The proof relies on the existence of the minimal perfect coloring for the bipartite representation of a hypergraph.

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A hypergraph  $\mathcal{G}$  is a covering of a hypergraph  $\mathcal{H}$ , if there exists a perfect coloring of  $\mathcal{G}$  whose parameter matrix  $\mathbb{S}$  is the full adjacency matrix of  $\mathcal{H}$ .

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**Equivalently:**  $\mathcal{G}$  is a *k*-covering of  $\mathcal{H}$  if one get the hypergraph  $\mathcal{H}$  uniting suitable groups of *k* vertices in  $\mathcal{G}$  and preserving the adjacency between them.

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**Equivalently:** G is a *k*-covering of H if one get the hypergraph H uniting suitable groups of *k* vertices in G and preserving the adjacency between them.

#### Theorem (T., 2021+)

If a hypergraph  $\mathcal{G}$  is a covering of  $\mathcal{H}$ , then every eigenvalue of  $\mathcal{H}$  is an eigenvalue of  $\mathcal{G}$ .

### Theorem (T., 2021+)

Let  $\mathcal{G}$  be a covering of a hypergraph  $\mathcal{H}$ . Then for every perfect coloring of  $\mathcal{H}$  with the parameter matrix  $\mathbb{S}$ , there is a perfect coloring of  $\mathcal{G}$  with the same parameter matrix  $\mathbb{S}$ .

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#### Theorem (Leighton, 1982)

Graphs  $H_1$  and  $H_2$  have the minimal perfect coloring with the same parameter matrix if and only if there exists a graph *G* covering both  $H_1$ and  $H_2$ .

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### Thank you for your attention!

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